Carnegie Mellon University

Database Query Optimization

Cascades Query Optimizer

ERRATA

Volcano Clarification \rightarrow Optimization Phases \rightarrow Enforcers

Send Corrections: <u>db-mistakes@cs.cmu.edu</u>

Start with a logical plan of what result the query needs to produce.

Generation Phase:

Logical Op

Physical Op

- → Apply transformation rules to generate all possible logical expression alternatives.
- **Cost Analysis Phase:**
- → Apply implementation rules to generate physical operators.



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Logical Op
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VOLCANO: ENFORCERS

Enforcers are physical operators that ensure the properties of the output of a sub-plan / expression.

Volcano's rule engine has additional logical to avoid considering operators below it in the plan that satisfy its property requirements. \rightarrow Example: INDEX_SCAN(xxx.b)



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Source: EQOP Book

TODAY'S AGENDA

Cascades Overview Tasks / Scheduling Optimizations Implementations Project #1

CASCADES OPTIMIZER

Object-oriented implementation of the previous Volcano query optimizer.

- → **Top-down approach** (backward chaining) using branchand-bound search.
- \rightarrow Depth-first search ordering of tasks via a stack.

Supports expression re-writing through a direct mapping function rather than an exhaustive search.



Graefe



EFFICIENCY IN THE COLUMBIA DATABASE QUERY OPTIMIZER PORTLAND STATE UNIVERSITY MS THESIS 1998

CASCADES: KEY IDEAS

Optimization tasks as data structures. \rightarrow Patterns to match + transformation rule to apply

Rules to place property enforcers.

 \rightarrow Ensures the optimizer generates correct plans.

Ordering of moves by promise.

 \rightarrow Dynamic task priorities to find optimal plan more quickly.

Unified representation of rules & operators.

 \rightarrow Single search engine for logical and physical operators.



CASCADES: EXPRESSIONS

An <u>expression</u> represents some operation in the query with zero or more input expressions.
→ Optimizer needs to quickly determine whether two expressions are equivalent.

```
SELECT * FROM A
JOIN B ON A.id = B.id
JOIN C ON C.id = A.id;
```

Logical Expression: (A \bowtie B) \bowtie C Physical Expression: (A_{Seq} \bowtie_{HJ} B_{Seq}) \bowtie_{NL} C_{Idx}

A **group** is a set of logically equivalent logical and physical expressions that produce the same output. \rightarrow All logical forms of an expression.

 \rightarrow All physical expressions derived from selecting allowable physical operators for the corresponding logical forms.

dno	Output: {ABC}	Logical Exprs 1. (A⋈B)⋈C 2. (B⋈C)⋈A	Physical Exprs1. $(A_{Seq} \bowtie_{SM} B_{Seq}) \bowtie_{SM} C_{Seq}$ 2. $(A_{Seq} \bowtie_{HJ} B_{Seq}) \bowtie_{HJ} C_{Seq}$
Gr	Properties: <i>None</i>	3. (A⋈C)⋈B 4. A⋈(B⋈C) ⋮	3. $(B_{Seq} \bowtie_{NL} C_{Seq}) \bowtie_{NL} A_{Seq}$ 4. $(A_{Seq} \bowtie_{NL} C_{Seq}) \bowtie_{NL} B_{Seq}$ 5.:

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Gr	Properties: <i>None</i>	3. (A⋈C)⋈B 4. A⋈(B⋈C) ⋮	3. $(B_{Seq} \bowtie_{NL} C_{Seq}) \bowtie_{NL} A_{Seq}$ 4. $(A_{Seq} \bowtie_{NL} C_{Seq}) \bowtie_{NL} B_{Seq}$ 5. :	

Equivalent Expressions

- A **group** is a set of logically equivalent logical and physical expressions that produce the same output. \rightarrow All logical forms of an expression.
- \rightarrow All physical expressions derived from selecting allowable physical operators for the corresponding logical forms.

Output: {ABC} Properties: None	Logical Exprs 1. (A⋈B)⋈C 2. (B⋈C)⋈A 3. (A⋈C)⋈B 4. A⋈(B⋈C)	Physical Exprs1. $(A_{Seq} \bowtie_{SM} B_{Seq}) \bowtie_{SM} C_{Seq}$ 2. $(A_{Seq} \bowtie_{HJ} B_{Seq}) \bowtie_{HJ} C_{Seq}$ 3. $(B_{Seq} \bowtie_{NL} C_{Seq}) \bowtie_{NL} A_{Seq}$ 4. $(A_{Seq} \bowtie_{NL} C_{Seq}) \bowtie_{NL} B_{Seq}$
None	:	5.:

Instead of explicitly instantiating all possible expressions in a group, the optimizer implicitly represents redundant expressions in a group with a placeholder (e.g., {ABC}).

→ This reduces the number of transformations, storage overhead, and repeated cost estimations.

Output:	Logical Exprs 1. {AB}⊠{C}	Physical Exprs 1. {AB}⊠ cu{C}
{ABC}	2. {BC}⋈{A} 3. {AC}⋈{B}	2. $\{AB\} \bowtie_{HJ} \{C\}$ 3. $\{AB\} \bowtie_{NL} \{C\}$
Properties: <i>None</i>	4. {A}⊠{BC} ∶	4. {BC}⊠ _{SM} {A} ∶

CASCADES: RULES

- A <u>**rule</u>** is a transformation of an expression to a logically equivalent expression.</u>
- → **Transformation Rule:** Logical to Logical
- → **Implementation Rule:** Logical to Physical

Each rule is represented as a pair of attributes:

- \rightarrow <u>**Pattern**</u>: Defines the structure of the logical expression that can be applied to the rule.
- \rightarrow **Substitute**: Defines the structure of the result after applying the rule.

CASCADES: RULES

Pattern



CASCADES: ENFORCERS

Represent enforcers as rules that change physical properties of a plan. → An enforcer inserts physical operators that change the physical properties of the plan.

Cascades does <u>not</u> require special casing to implement enforcers. \rightarrow Volcano integrated enforcers as operators

→ volcano integrated enforcers as operators via ad hoc code.



CASCADES: TASKS

- A <u>task</u> is a fine-grained unit of work that represents an operation in the query optimization process.
- → Break down optimization into smaller, manageable pieces to allow for more flexible and efficient exploration.

The order that a task is invoked can be customized via **promise**.

CASCADES: PROMISES

A task's **promise** is the estimated benefit of a move on a given expression relative to other tasks.

→ Example: Give a higher priority to a join order swap rule if the outer table expression is larger than inner table expression.

Optimizer must still ensure tasks execute in the order defined by their dependencies.



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CASCADES: TASKS

#1 – Optimize Group:

 \rightarrow Generate best physical plan for a group.

#2 – Optimize Expression:

 \rightarrow Generate best physical plan for a specific expression.

#3 – Explore Group:

 \rightarrow Generate logical expressions for a group.

#4 – Explore Expression:

 \rightarrow Generate logical transformations for a specific expression.

#5 – Apply Rule:

 \rightarrow Apply a rule to an input expression.

#6 – Optimize Inputs:

 \rightarrow Optimize the inputs of a given expression.



CASCADES: TASK-BASED SEARCH

Optimizer maintains a LIFO stack of tasks to perform actions on groups and expressions.

Stack ensures expressions are derived after the best plans of its input expressions are derived.

- \rightarrow Removes sequential ordering of independent optimizations.
- \rightarrow Tasks are stored in the heap rather than in the program stack to reduce OOM errors.

Stores all previously explored alternatives in a compact graph structure / hash table.

Equivalent operator trees and their corresponding plans are stored together in groups.

Provides an overview of the optimizer's search progress that is used in multiple ways:

- → Transformation Result Memorization
- \rightarrow Duplicate Group Detection
- \rightarrow Property + Cost Management
- \rightarrow Superfluous Rule Evaluation

PRINCIPLE OF OPTIMALITY

Every sub-plan of an optimal plan is itself optimal.

This allows the optimizer to restrict the search space to a smaller set of expressions.

→ The optimizer never has to consider a plan containing subplan P1 that has a greater cost than equivalent plan P2 with the same physical properties.



	Best Expr	Cost	
{ABC}			
{AB}			
{A}			
{B}			
{C}			

Outpu	it:	Logical Exprs	Physical Exprs
{ABC	}	1. {AB}⋈{C}	
Properti <i>None</i>	ies:		





















Cost: 10

Cost: 20

Output:	Logical Exprs	Physical Exprs	Output	Logical Exprs	Phy	sical Exprs
	1. GET(A)	1. SeqScan(A)		1. GET(B)	1.	SeqScan(B)
143		2. IdxScan(A)	103		2.	IdxScan(B)
Properties: <i>None</i>			Properties: <i>None</i>			





Cost: 10

Cost:	20
-------	----

Output	Logical Exprs	Physical Exprs	Output	Logical Exprs	Phy	sical Exprs
	1. GET(A)	 SeqScan(A) 		1. GET(B)	1.	SeqScan(B)
{A}		2. IdxScan(A)	{D}		2.	IdxScan(B)
Properties: <i>None</i>			Properties: <i>None</i>			





Cost: 10

Cost: 2	20
---------	----

Output	Logical Exprs	Physical Exprs	Output	Logical Exprs	Phy	sical Exprs
	1. GET(A)	 SeqScan(A) 	(P)	1. GET(B)	1.	SeqScan(B)
{A}		2. IdxScan(A)	{D}		2.	IdxScan(B)
Properties: <i>None</i>			Properties: <i>None</i>			



Cost: 10

Cost: 20

						-		
Output:	Logical Exprs	Physical Exprs		Output	Logical Exprs	Physical Exprs		
	1. GET(A)	 SeqScan(A) 			1. GET(B)	1. SeqScan(B)		
{A}		2. IdxScan(A)		{D}		2. IdxScan(B)		
Properties: <i>None</i>				Properties: <i>None</i>				

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Cost: 10

Cost: 20

Output: {A}	Logical Exprs	Physical Exprs		Output: Logical Exprs		Physical Exprs	
	1. GET(A)	 SeqScan(A) 		{B}	1. GET(B)	1.	SeqScan(B)
		2. IdxScan(A)				2.	IdxScan(B)
Properties: <i>None</i>				Properties: <i>None</i>			

	Best Expr	Cost
{ABC}		
{AB}	80	
{A}	SeqScan(A)	10
{B}	SeqScan(B)	20
{C}	IdxScan(C)	5

Output: {ABC} Properties: None

Cost: 50+(10+20)



Output: {AB}	Logical Exprs 1. {A}⋈{B} 2. {B}⋈{A}	Physical Exprs 1. $\{A\} \bowtie_{M} \{B\}$ 2. $\{A\} \bowtie_{u_{1}} \{B\}$	Output: {C}	Logical Exprs 1. GET(C)	Physical Exprs 1. SeqScan(C) 2. IdxScan(C)
Properties: <i>None</i>		3. $\{B\} \bowtie_{NL} \{A\}$ 4. $\{B\} \bowtie_{HJ} \{A\}$	Properties: <i>None</i>		

Cost: 20

Cost: 10

Output: {A}	Logical Exprs	Physical Exprs		Output: Logical Exprs		Physical Exprs		
	1. GET(A)	 SeqScan(A))	{B}	1. GET(B)	1. SeqScan(B)		
		2. IdxScan(A)				2. IdxScan(B)		
Properties: <i>None</i>				Properties: <i>None</i>				

	Best Expr	Cost	
{ABC}	$({A}M_{HJ}{B})M_{HJ}{C}$	125	
{AB}	$\{A\}\bowtie_{HJ}\{B\}$	80	
{A}	SeqScan(A)	10	
{B}	SeqScan(B)	20	
{C}	IdxScan(C)	5	

Cost: 40+(80+5)

Output	Logical Exprs	Physical Exprs			
	1. {AB}⊠{C}	1. {AB}⊠ _{NL} C			
{ADC}	2. {BC}⊠{A}	2. {BC}⊠ _{NL} A			
Properties:	3. {AC}⋈{B}	3. {AC}⊠ _{NL} B			
None	4. {B}⊠{AC}	:			

Cost: 50+(10+20)

Cost: 5

Output: {AB}	Logical Exprs 1. {A}⋈{B} 2. {B}⋈{A}	Physical Exprs 1. {A}⋈ _N {B} 2. {A}⋈ _{HI} {B}	Output: {C}	Logical Exprs 1. GET(C)	Physical Exprs 1. SeqScan(C) 2. IdxScan(C)
Properties: <i>None</i>		3. $\{B\} \bowtie_{NL} \{A\}$ 4. $\{B\} \bowtie_{HJ} \{A\}$	Properties: <i>None</i>		

Cost: 20

Cost: 10

Output: {A}	Logical Exprs	Physical Exprs	Output, Log		Logical Exprs	s Physical Ex	
	1. GET(A)	 SeqScan(A) 	out		1. GET(B)	1.	SeqScan(B)
		2. IdxScan(A)		{D}		2.	IdxScan(B)
Properties: <i>None</i>				Properties: <i>None</i>			

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OBSERVATION

Promises enable Cascades to potentially target more beneficial transformations more quickly in the search process.

Cascades' flexible architecture enables other optimizations to reduce search times.

SIMPLIFICATION RULES

Some rules simplify the logical query plan and almost always reduce its cost.

Instead of creating alternative expressions, a **<u>simplification rule</u>** replaces the expression with the transformed expression.

- \rightarrow Removes the need to retain unnecessary state.
- \rightarrow Equivalent to Starburst rewrite rules.

MACRO RULES

To reduce the complexity of the search space and find a better plan more quickly, the optimizer can support **macro rules** that apply multiple transformations in a single rule.

This goes against the spirit of simple, independent rules as defined in the Volcano paper.

Source: EQOP Book

PARALLEL SEARCH

If tasks are independent, then the optimizer can execute them in parallel on multiple threads. \rightarrow Memo table is the shared state of the optimization process.

 \rightarrow Need to ensure internal data structures are thread-safe.

AFAIK, <u>Orca</u> is the only multi-threaded Cascades optimizer implementation.

We will discuss this more in <u>Lecture #17</u>.

Source: EQOP Book

CASCADES IMPLEMENTATIONS

Standalone:

- \rightarrow <u>Wisconsin OPT++</u> (1990s)
- \rightarrow <u>Portland State Columbia</u> (1990s)
- \rightarrow <u>Greenplum Orca</u> (2010s)
- \rightarrow <u>CMU optd</u> (2025)

Integrated:

- \rightarrow Microsoft SQL Server (1990s)
- \rightarrow <u>Tandem NonStop SQL</u> (1990s)
- $\rightarrow \underline{\text{Clustrix}}$ (2000s)
- \rightarrow <u>CockroachDB</u> (2010s)
- \rightarrow <u>Snowflake</u> (2010s)
- \rightarrow <u>Databricks</u> (2010s)

MICROSOFT SQL SERVER

First Cascades implementation started in 1995.

- \rightarrow Derivatives are used in many MSFT database products.
- \rightarrow All transformations are written in C++. No DSL.
- \rightarrow Scalar / expression transformations are written in procedural code and not rules.

DBMS applies transformations in multiple stages with increasing scope and complexity.

→ The goal is to leverage domain knowledge to apply transformations that you always want to do first to reduce the search space.

MSSQL: MULTI-STAGE OPTIMIZATION



Source: Nico Bruno + Cesar Galindo-Legaria

MSSQL: OPTIMIZATIONS

Optimization #1: Timeouts are based on the number of transformations not wallclock time.
 → Ensures that overloaded systems do not generate different plans than under normal operations.

Optimization #2: Pre-populate the Memo Table with potentially useful join orderings.
→ Heuristics that consider relationships between tables.
→ Syntactic appearance in query.

GREENPLUM ORCA

Standalone Cascades implementation in C++.

- \rightarrow Originally written for <u>Greenplum</u>.
- \rightarrow Extended to support <u>HAWQ</u>.
- \rightarrow Supports multi-threaded search.

A DBMS integrates Orca by implementing API to send catalog + stats + logical plans and then retrieve physical plans.

Open-sourced in 2010s but then closed-sourced in 2024 after Broadcom acquisition.



COCKROACHDB

Custom Cascades implementation written in 2018.
All transformation rules are written in a custom
DSL (<u>OptGen</u>) and then codegen into Go-lang.
→ Can embed Go logic in rule to perform more complex analysis and modifications.

Also considers scalar expression (predicates) transformations together with relational operators.

Source: <u>Rebecca Taft</u>

COCKROACHDB

Custom Cascades im All transformation r DSL (<u>OptGen</u>) and t → Can embed Go logic i analysis and modifica

Also considers scalar transformations toge

DSL: Optgen

// ConstructNot constructs an expression for the Not operator.
func (_f *Factory) ConstructNot(kinput opt.ScalarExpr) opt.ScalarExpr {

```
// [EliminateNot]
{
    _not, _ := input.(*memo.NotExpr)
    if _not != nil {
        input := _not.Input
        if _f.matchedRule == nil || _f.matchedRule(opt.EliminateNot) {
            _expr := input
            return _expr
        }
    }
    // ... other rules ...
    e := _f.mem.MemoizeNot(input)
    return _f.onConstructScalar(e)
}
Cockroach LABS
```

Source: <u>Rebecca Taft</u>

PARTING THOUGHTS

Andy prefers the Cascades' architecture because it is conceptually cleaner than other implementations.
→ The quality of a plan is still highly dependent on accurate statistics and cost estimations.

The Germans disagree with top-down search...

PROJECT #1 - QUERY OPTIMIZER EVALUATION

You will use <u>Apache Calcite</u> to optimize SQL queries that are then executed on Calcite (via the <u>Enumerable</u> adapter) and <u>DuckDB</u>.



Due Date: Friday Feb 28th

https://15799.courses.cs.cmu.edu/spring2025/project1.html

PROJECT #1 - TASKS

We provide starter code to build and run a Calcite application. You need to implement the highlighted region.



Graded with: optimize.sh workload.tgz output_dir

For each foo.sql in workload.tgz/queries, write to output_dir the following files:

foo.sql foo.txt foo_results.csv foo_optimized.sql

[Original SQL query] [Logical Calcite Plan] foo_optimized.txt [Enumerable Calcite Plan] [Results from foo_optimized] [Optimized SQL query]

You are graded against the reference solution on:

- RelRunner result correctness
- DuckDB performance with disable_optimizer

DEVELOPMENT HINTS

Follow the roadmap.

Use our code for serializing the txt and csv files. Beware of case sensitivity issues (e.g., table name).

The "<u>standard library</u>" of rules does <u>**not**</u> include operation-specific rules. You may need them.

If there's limited support for optimizing or deparsing Foo, maybe you can turn Foo into Bar. Then optimize and/or deparse that instead.

THINGS TO NOTE

Do <u>**not**</u> hardcode based on the provided workload's data distribution – we test on a different workload.

Limit your memory usage to avoid autograder crashes: java -Xmx4096m jar App.jar APP_ARGS

Post your questions on Piazza.

Make sure you submit a reflection that documents what you tried. Bullet points are fine.

NEXT CLASS

Cascades Transformations

- \rightarrow Everything but sub-queries.
- \rightarrow You do <u>not</u> need to submit a reading synopsis.