Special Topics:

Self-Driving Database Management Systems

Partitioning

@Karthik R // 15-799 // Spring 2022

LAST CLASS

- Design Choices in Knob Tuning
 - Knob Selection
 - Configuration Optimization
 - Knowledge Transfer
- Survey of Algorithms



TODAY'S AGENDA

- Refresher on Partitioning
- What is Microsoft up to?
- The Partitioning Problem
- Approach
- Results
- Parting Thoughts



TODAY'S AGENDA

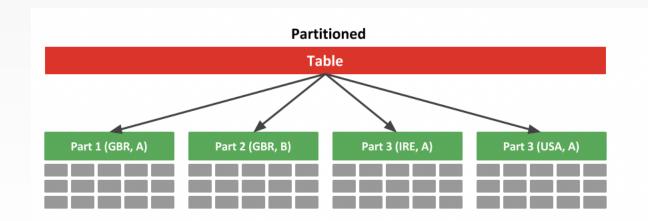


- Refresher on Partitioning
 - WHY, WHAT and HOW of Partitioning
 - The Big Picture
- What is Microsoft up to?
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WHAT IS PARTITIONING?

Partitioning splits database across multiple resources





NEED FOR PARTITIONING

When the server crashes and you realize back-up was on that same server



NEED FOR PARTITIONING

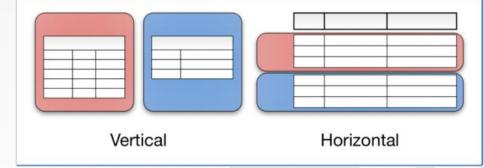
Improves:

- Scalability
 - → "Scale out" rather than "Scale up"
- Performance
 - → Queries run over subset of data
- Availability
 - → Removes single point of failure
 - → Related: Concept of "replicas" or redundancy

TYPES OF PARTITIONING

- By design
 - Horizontal, Vertical, ...

- By method
 - Hash, Range, List, ...



- By Subject
 - Tables, (n-Cl) Indexes, Materialized Views, Partitions

THE BIG PICTURE

- Physical Database Design involves
 - → Indexes, Materialized Views and <u>Partitions</u>

- Performance vs Manageability (for humans)
- Automated partition selection is NP-hard
 - → Combinatorial explosion
 - → Superficially similar to index selection

TODAY'S AGENDA



- Refresher on Partitioning
- What is Microsoft up to?
 - Timeline
 - Database Tuning Advisor
- The Partitioning Problem
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MICROSOFT'S TIMELINE

1997-98

- Cost-driver Index Selection
- Index Analysis

2000

Selection of Materialized Views

2004

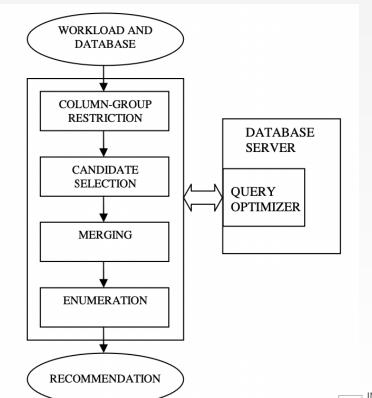
- Release of DB Tuning Advisor (DTA) for MSSQL 2005
- Integration of horizontal and vertical partitioning

DATABASE TUNING ADVISOR

Considerations

- → Co-location of Join columns (size and locality)
- → Mutual exclusivity (unlike indexes)
- → Specificity vs Generality
- → Storage and Update costs
- → Alignment

DATABASE TUNING ADVISOR



DATABASE TUNING ADVISOR

- Column Group Restriction
 - → Simulating vertical partitions is hard!
 - → Heuristic based
- Candidate Selection
 - → Greedy (m, k) algorithm
- Merging
 - → Over the entire workload
 - → Involves indexes, MVs and partitions
- Enumeration



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PARTITIONING FOR CLOUD DBs

"Learn the optimal partitioning for each cloud customer, for a given database schema, for a given workload"

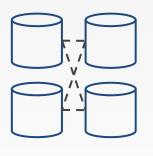
Suitability of Reinforcement Learning

- → Combinatorial Optimization problem
- → Exploration vs Exploitation instead of Greedy

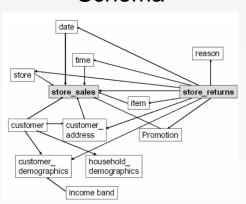


PARTITIONING FOR CLOUD DBs

Diverse Hardware



+ Complex Schema



+ Arbitrary Workload

```
AS ctr customer sk,
SELECT sr customer sk
                          AS ctr store sk,
       sr store sk
      Sum(sr return amt) AS ctr total return
FROM
      store returns,
      date dim
WHERE sr returned date sk = d date sk
      AND d vear = 2001
CDOLLD By or customer ch
```

(+ Various DBMS)



Postgres-XL

- Problems with existing approaches
 - → Cost estimates are coupled to hardware
 - → Cost estimates themselves are inaccurate



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• State:

- \rightarrow Given a table T_i with attributes $(a_{i1}, a_{i2}, \dots, a_{in})$
- \rightarrow One-hot encoding as $(r_i, a_{i1}, a_{i2}, \dots, a_{in})$
- → "Edges" for co-partitioning
- → Workload state
- → All appended together



- Workload state modeling:
 - → Encoding JOIN predicates, WHERE clauses etc

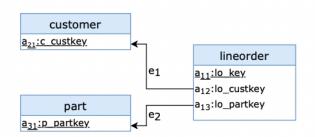
 Does not account for arbitrary nesting.
 - → Nested query featurization complex encoding, larger training data
 - → Keep It Simple Stupid! Encode only frequency info. $s(Q) = (f_1, f_2, ..., f_m)$
 - → Bucketize queries based on Selectivity. How?

• Action:

- → Q-learning ... small state space desirable
- → Partitioning OR Replication
- → Only considers <u>HASH Partitions</u>
- → Only considers horizontal partitioning with fixed no. of nodes
- → One-hot encoding of (repl, partition, (de)activation) with at most one active







q₁: SELECT * FROM customer c, lineorder l WHERE I.lo_custkey=c.c_custkey;

q₂: SELECT * FROM part p, lineorder l WHERE I.lo_partkey=p.p_partkey;

(a) Database and Workload

Foreign-Key Edges:

Edge e_1 for lo_custkey \rightarrow c_custkey: active Edge e_2 for lo_partkey \rightarrow c_partkey: inactive $s(E) = (e_1, e_2) = (1, 0)$

Table States:

lineorder partitioned by lo_custkey $s(lineorder) = (r_1, a_{11}, a_{12}, a_{13}) = (0, 0, 1, 0)$

customer partitioned by c_custkey $s(customer) = (r_2, a_{21}) = (0, 1)$

part replicated $s(part) = (r_3, a_{31}) = (1, 0)$

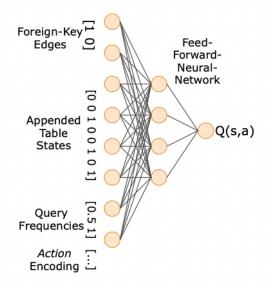
Query Frequencies:

 q_2 occurs twice as frequently as q_1 $s(Q) = (f_1, f_2) = (0.5, 1)$

(b) State Representation

(c) Q-Network with Encoded State

Figure 2: State Representation of Simplified SSB Schema and Workload.



• Cost:

- → "Network-centric" cost model for offline training
- → Runtime of queries for online training

• Reward:

- → Gain in performance for a workload
- $\rightarrow \mathbf{r} = -\sum_{j=1}^{m} f_j * c(P, q_j)$
- → Excludes cost of repartitioning for OLAP Workloads



OFFLINE TRAINING

- Enumeration join orderings
- Estimate optimal join strategy for each join using cost model.
- Sum of costs (network + compute) is the cost of a query
- Each iteration ('episode') comprises of $(t_{MAX} >= |T|)$ actions
- Train Q-network with SGD and loss



ONLINE TRAINING

- Runs on a copy of the database, workloads
- Cost model = true runtime
- Enumeration-based training \rightarrow v.v. expensive!
- Optimization 1: Sampling
 - Use of a scale factor (per query) to weigh the costs
 - Very small samples can lead to suboptimal partitions
 - What should the rate/threshold be?



ONLINE TRAINING

- Optimization 2: Query Runtime Caching
 - Maintain a cache of query runtimes per partition
 - Given two states s_a , s_b , their partitions P_a , P_b , a query q_i
 - We need to estimate costs, only if q_i queries $t \in |P_b P_a|$
- Optimization 3: Lazy Repartitioning
 - Repartitioning take time!
 - We partition only if q_i queries $t \in |P_b P_a|$
- Optimization 4: Timeouts
 - If a query costs more in P_b, it is not a good partition!



WORKLOAD CHANGES

• Committee of Experts:

- Different queries favor different optimal partitions
- Poison the freq. vector to extract these "reference partitions"
- A query belongs to subspace of a reference partition if $P_i = argmax_{P_i \in \{P_1...Pn\}} \sum_{j=1}^m f_j S_j * c_{sample}(P, q_j)$
- Train an agent for each subspace

• Incremental Training:

- New query frequencies added to input state
- Runtime Cache speeds things up
- Might not need to train new reference partition

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WORKLOADS AND SETUP

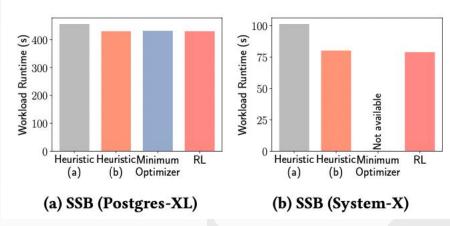
- Analytical Workloads:
 - → Star Schema Benchmark
 - \rightarrow TPC-DS
 - → TPC-CH
- Setup:
 - → Postgres-XL (open source, disk-based)
 - → System-X (commercial, memory-based)
 - \rightarrow 4-6 node clusters
 - → 128GB RAM, 2*10-core Intel Xeon CPUs





BASELINE

- Compared to:
 - Heuristic A
 - Heuristic B
 - Minimum Optimizer



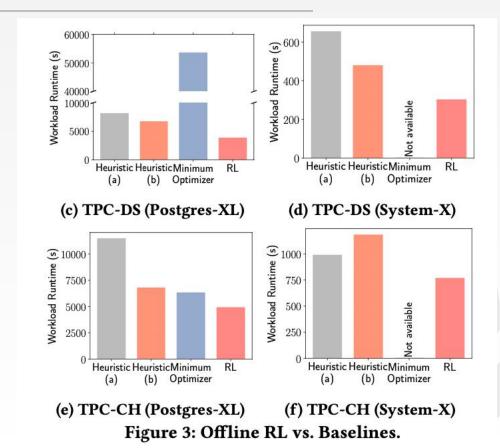
- Heuristic A: Most frequently joined dim. table
- Heuristic B: Largest table
- Minimum optimizer: Non ML opt. algorithm





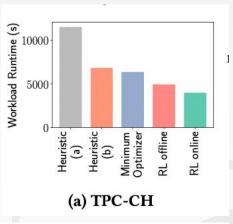
OFFLINE TRAINING RESULTS

- For TPC-DS, DRL Agent suggests superior partitions
- For TPC-CH, DRL Agent optimizes for network costs, leading to the difference in runtimes



ONLINE TRAINING RESULTS

• Online training has better results because of a difference in cost model



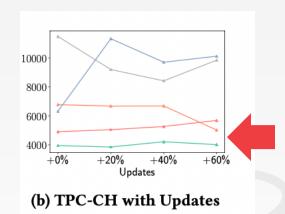
Optimizations	Training Time	Speedup
None	4621h	-
+ Runtime Cache	1160.4h	4.0
+ Lazy Repartitioning	60h	19.3
+ Timeouts	33.4h	1.8
+ Offline Phase	13.3h	2.5

Table 2: Training Time Reduction of Optimizations.



ADAPTIVITY TO CHANGE

- New data Advisor performs fine for "non-significant" changes to data.
 Requires re-training otherwise.
- Changing Workload Mix Committee of experts performs well.
- New queries Requires incremental training; bootstrapped agent with low exploration.



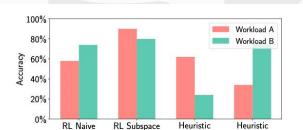


Figure 5: Best Partitioning found by Different Approaches for Varying Workloads (higher is better).

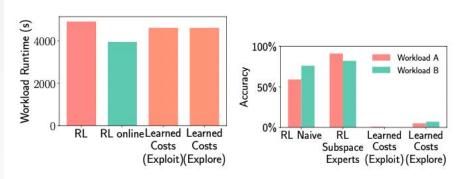
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OTHER ML APPROACHES: COMPARISON

Compared with NEO, a learned query optimizer

- Exp 1: Static Workload
- Exp 2: Workload Adaptivity



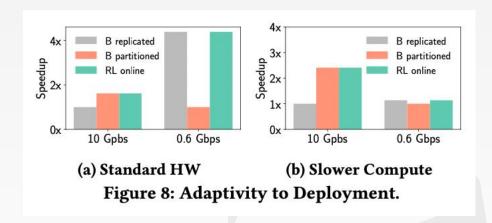
(a) TPC-CH Schema Figure 7: RL vs. Neural Baselines.

(b) Workload Adaptivity



ADAPTIVITY TO DEPLOYMENT

- Run on System X, to avoid disk access costs
- Compares trade-off between Compute (colocation) and Network (parallelism)



PARTING THOUGHTS

- Subset of the Partitioning Problem
 - → Only Horizontal, Hash Partitions
 - → Only evaluated on OLAP workloads
- Choice of Network Model
- Susceptible to Schema Changes?
- Accuracy as a Heuristic in Experiments
- Heterogenous Architecture/Hardware



NEXT CLASS

• Lin Ma's paper on Workload Modeling

